

# Cluster Algebras and Binary Words

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## Binary word

A *binary word*  $w = w_1w_2 \dots w_d$  is a finite (possibly empty) sequence of letters on the alphabet  $\{0, 1\}$  starting with 1. A *subword*  $s$  of a binary word is a subsequence of  $w$  which itself is a binary word.

- Each non-empty binary word can be associated with (the Hasse diagram of) a piece-wise linear poset  $P$ .

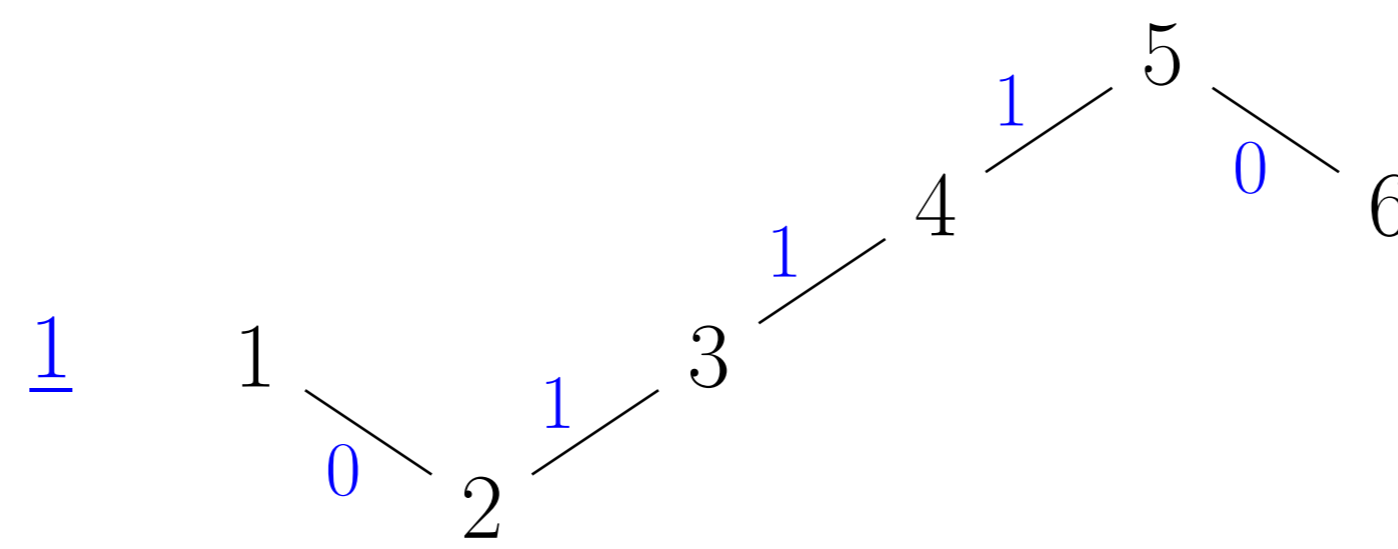


Figure 1. The Hasse diagram associated to the word 101110.

## Antichain

An *antichain* is a subset  $A = \{A_1, A_2, \dots, A_r\}$  of a poset such that no two distinct elements in  $A$  are comparable.

Example:  $\{1, 3, 6\}$ ,  $\{1, 4\}$  and  $\{2, 6\}$  are antichains of the poset corresponding to the above Hasse diagram, while  $\{2, 4\}$  is not.

## Trie of subwords (Leroy, Rigo, and Stipulanti)

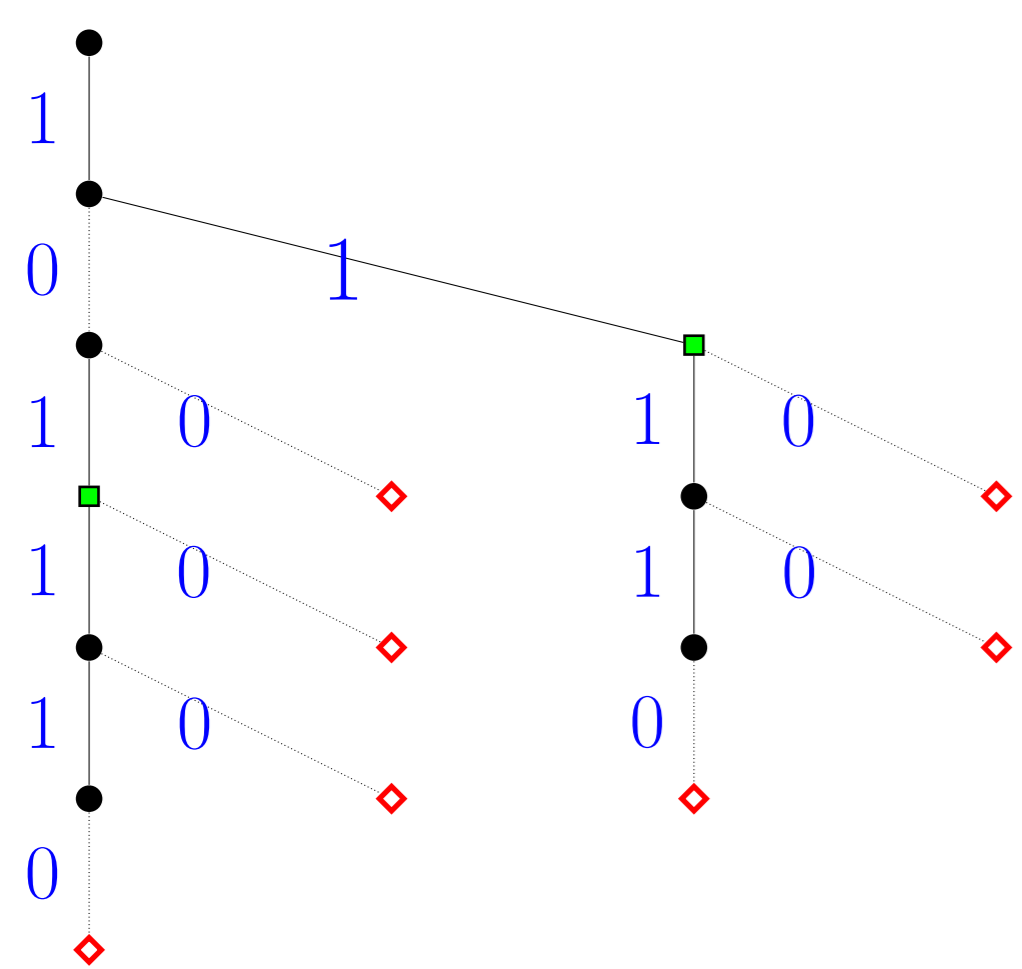


Figure 2. Trie of subwords corresponding to  $w = 101110$ .

## Antichain trie

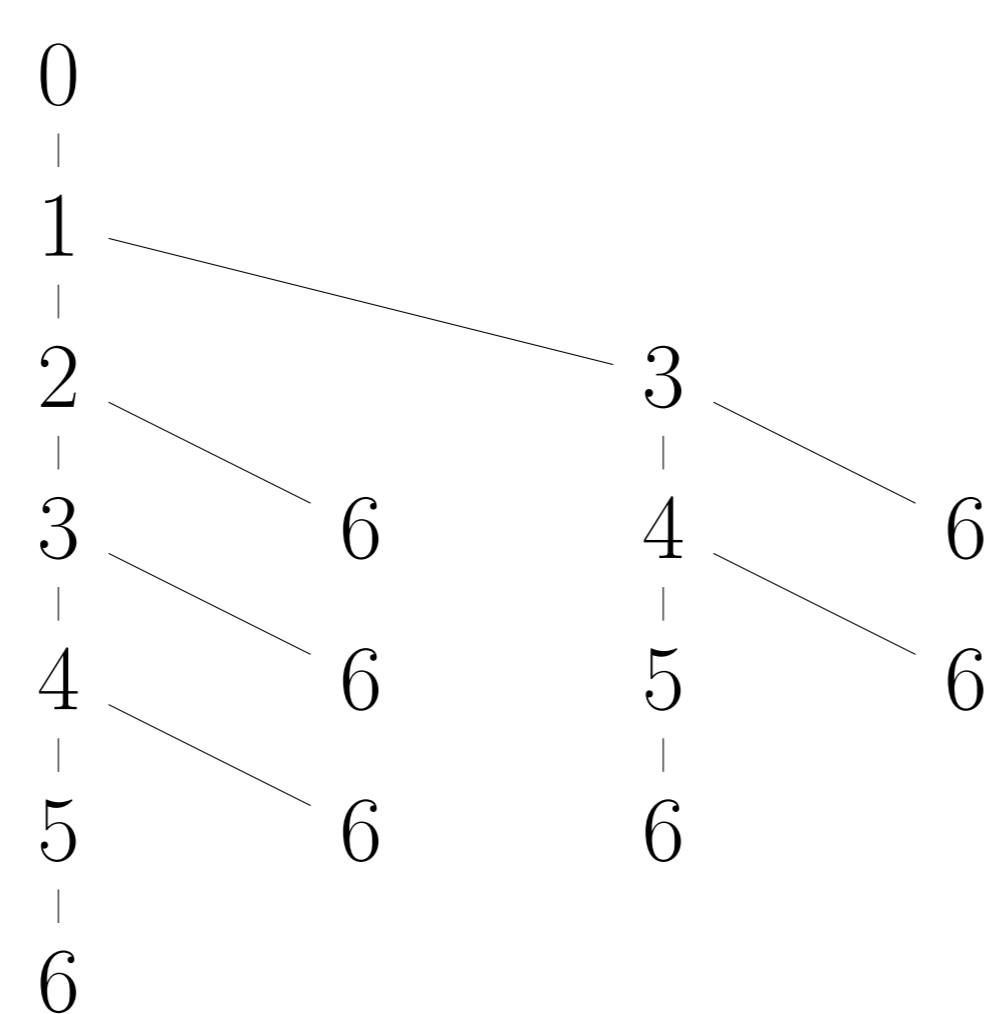


Figure 3. Antichain trie corresponding to  $w = 101110$ .

## Corresponding antichains

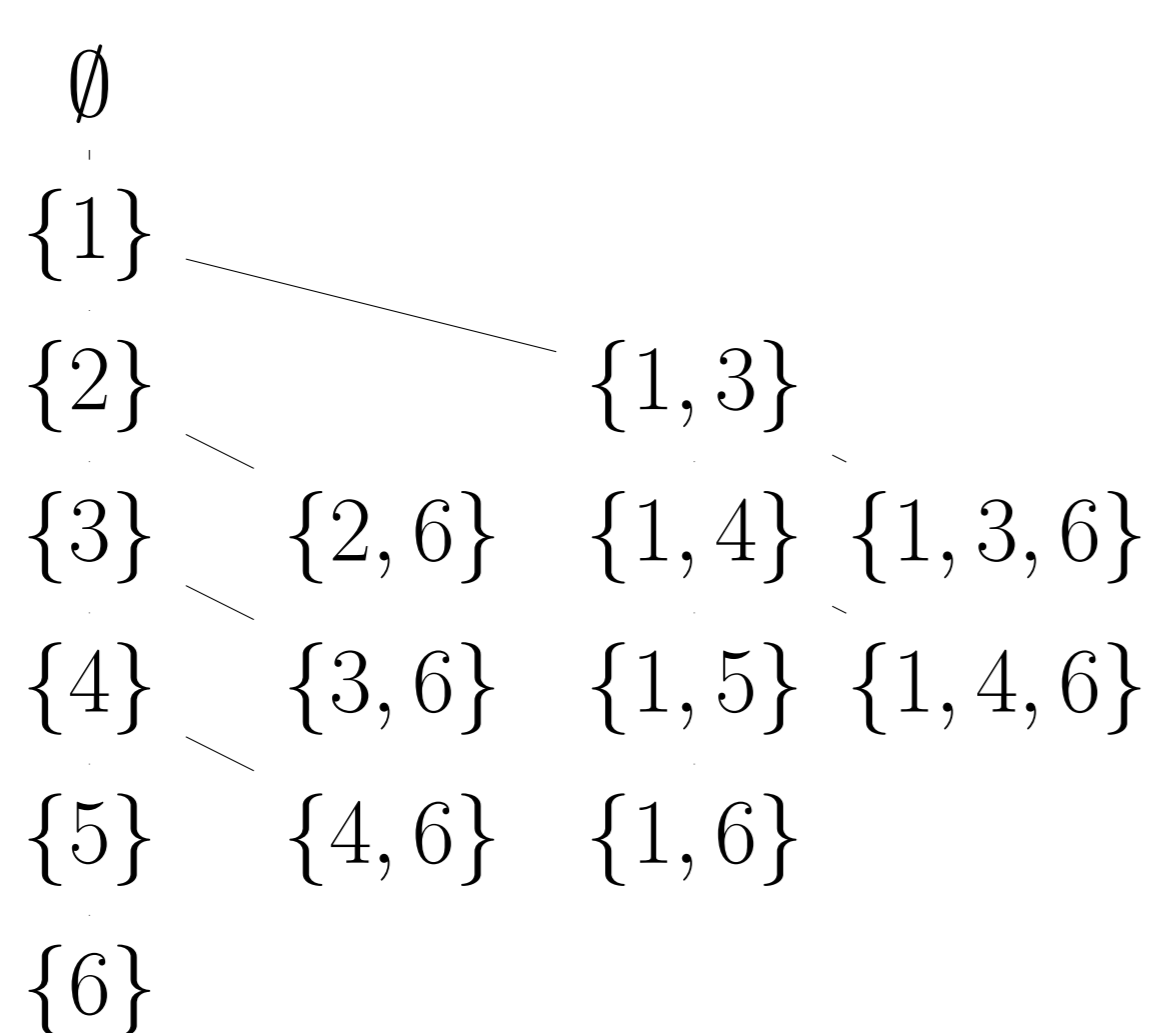
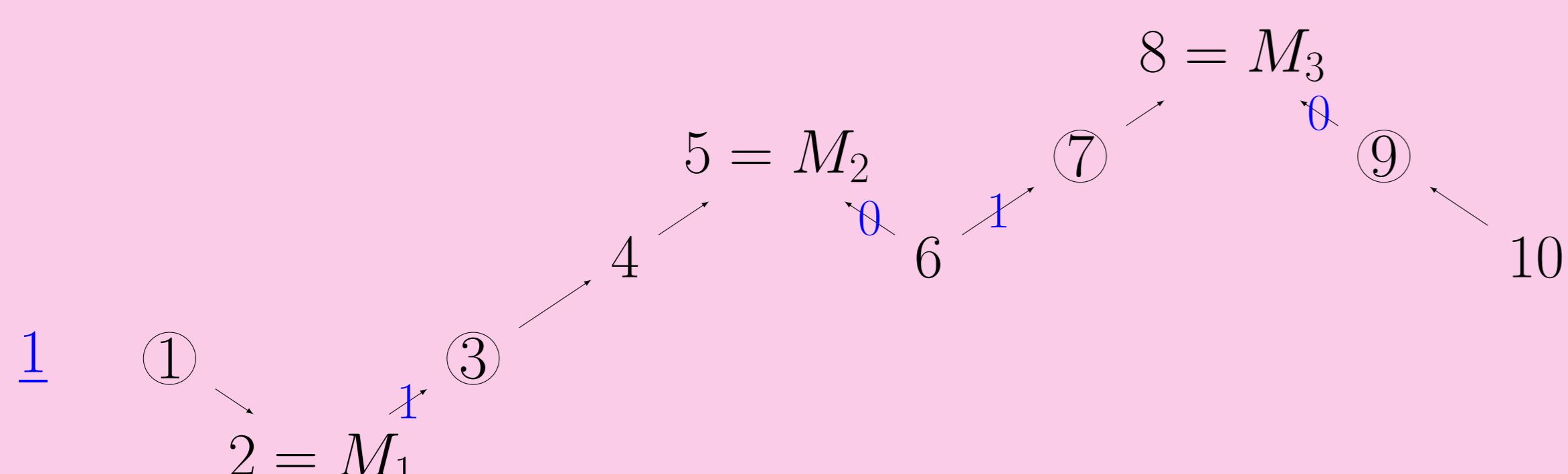


Figure 4. Trie of antichains for  $w = 101110$ .

## Definition

We define a map from the antichains in  $P$  to the subwords of  $w$ . For example, the antichain  $A = \{A_1 = \{1\}, A_2 = \{3\}, A_3 = \{7\}, A_4 = \{9\}\}$  is mapped to the subword  $s = \underline{1}1010$ .



## Theorem

The map given above is a bijection from the antichains in  $P$  to the subwords of  $w$ .

## Snake graph

A *snake graph*  $G$  is a nonempty connected sequence of  $d$  square tiles constructed so that each new tile is glued to the north or east side of the previous tile.

- A *perfect matching* of  $G$  is a subset of edges of  $G$  such that every vertex of  $G$  is adjacent to exactly one edge. The *minimal matching*  $P_{min}$  is the unique perfect matching of  $G$  which contains the first south edge and only boundary edges.

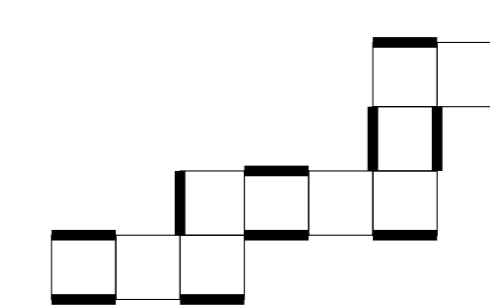


Figure 5. The minimal matching for  $w = 1011101100$ .

## From binary words to snake graphs

- A *cluster algebra* is a commutative algebra with distinguished generators called *cluster variables* which can be written as Laurent polynomials with positive coefficients. Certain Laurent polynomials can be associated to snake graphs.
- A *sign function* on a snake graph  $G$  is a map from the set of edges of  $G$  to  $\{0, 1\}$ .
- We can associate to each binary word  $w = 1w_2 \dots w_d$  a snake graph  $G(w)$ .

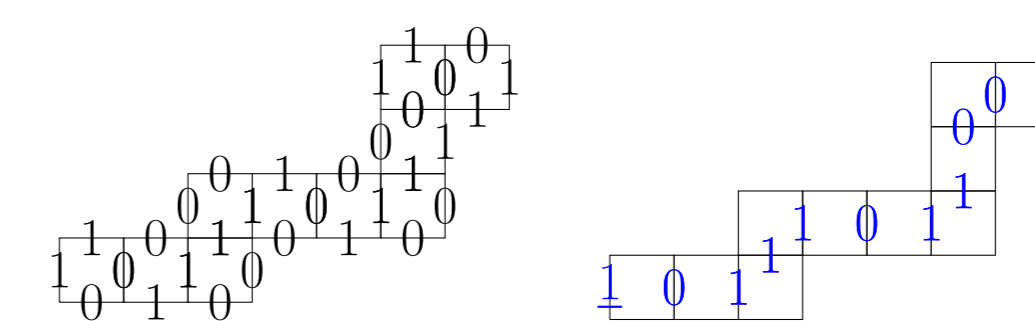


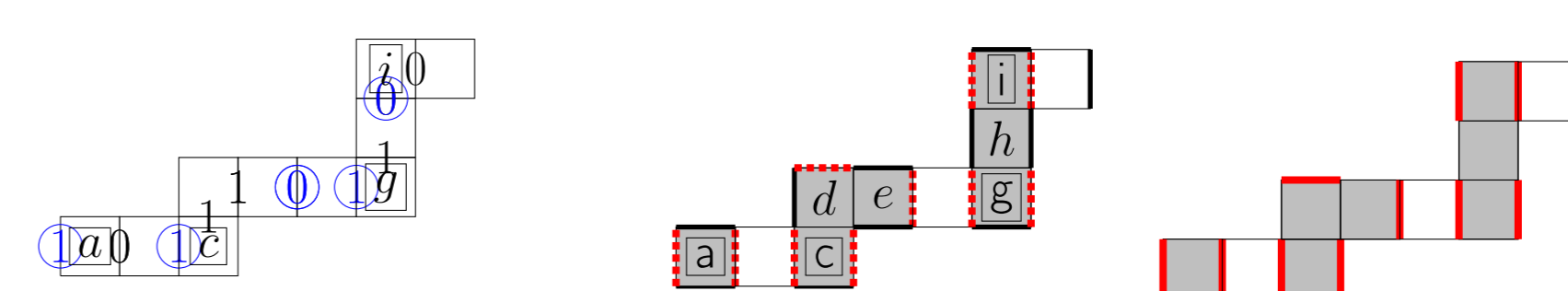
Figure 6. The sign function on the snake graph associated to the word  $w = 1011101100$ .

## Theorem

Given a binary subword  $s$  and its corresponding snake graph  $G = G(w)$ , the following map  $pm$  from the subwords of  $w$  to the perfect matchings of  $G$  is a bijection:

- Let  $s$  be a subword of  $w$ . If  $s$  is the empty word, let  $pm(s)$  be  $P_{min}$ . Otherwise, write  $s = 1w_{i_2} \dots w_{i_k}$  in such a way that each index  $i_k$  is as small as possible and circle the edges of  $G$  corresponding to the sign sequence for  $s$ .
- For each block  $L$  of consecutive circled edges, let  $\square_L$  be the tile which is immediately north/east of the last edge in  $L$ .
- Let  $fil(\square_L)$  be the smallest connected sequence of tiles containing  $\square_L$  such that the set of edges bounding  $fil(\square_L)$  not in  $P_{min}$  forms a perfect matching of  $fil(\square_L)$ .
- Let  $fil(s) = \bigcup_L fil(\square_L)$ , and define  $pm(s)$  to be the symmetric difference  $\{\text{edges bounding } fil(s)\} \ominus P_{min}$ .

## Example



The tiles  $\square_L$  associated to blocks  $L$  of circled edges (left); the set  $fil(s)$  of shaded tiles (center) and the set  $pm(s)$  of thick solid edges (right) for the subword  $s = \underline{1}1010$  of  $w = 1011101100$ .